

Math 210C Lecture 18 Notes

Daniel Raban

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1 Derived Functors and Injective Resolutions

1.1 Left derived functors

Let \mathcal{C} be an abelian category. Let $F : \mathcal{C} \rightarrow \mathcal{D}$ be an additive functor of abelian categories and let $f : P \rightarrow Q$ in \mathcal{C} . Then $F(f_*) : H_i(F(P)) \rightarrow H_i(F(Q))$ is “well-defined” in \mathcal{D} .

Proposition 1.1. *Let $F : \mathcal{C} \rightarrow \mathcal{D}$ be an additive functor. For each $A \in \mathcal{C}$, choose a projective resolution $P^\bullet \rightarrow A$. For each $A : A \rightarrow B$ choose $f : P^\bullet \rightarrow P^\bullet$ augmenting g . Then $L_i F : \mathcal{C} \rightarrow \mathcal{D}$ given by $L_i F(A) = H_i(F(P^\bullet))$ and $L_i F(g) = F(f_*)$ is a functor (unique up to unique isomorphism).*

Definition 1.1. $L_i F$ is called the i -th left derived functor of F .

Proof. This mostly follows from a proposition from last time. Check the remaining properties, such as composition. \square

Lemma 1.1. *If F is right exact, then $L_0 F \cong F$.*

Proof. Let $P_\bullet \rightarrow A \in \mathcal{C}$ in a projective resolution of A . Then $L_i F(A) = H_i F(P_\bullet)$, so $L_0 F(A) = H_0(F(P_\bullet))$, so

$$P_1 \longrightarrow P_0 \xrightarrow{\varepsilon} A \longrightarrow 0$$

gives

$$F(P_1) \xrightarrow{F(d_1)} F(P_0) \xrightarrow{\varepsilon} F(A) \longrightarrow 0$$

Then

$$\begin{array}{ccc} H_0(F(P_\bullet)) = \operatorname{coker}(F(d_1)) & \xrightarrow{\cong} & F(A) \\ \downarrow F(f)_* & & \downarrow F(g) \\ H_0(F(Q_\bullet)) & \xrightarrow{\cong} & F(B) \end{array}$$

and the left hand side of this diagram is actually

$$\begin{array}{c} L_0F(A) \\ \downarrow L_0F(g) \\ L_0F(B) \end{array}$$

□

Theorem 1.1. *Let $F : \mathcal{C} \rightarrow \mathcal{D}$ be right exact. For each short exact sequence*

$$0 \longrightarrow A \longrightarrow B \longrightarrow C \longrightarrow 0$$

in \mathcal{C} , there exists a long exact sequence

$$\begin{array}{ccccccc} \cdots & \longrightarrow & L_2F(C) & \xrightarrow{\delta} & L_1F(A) & \longrightarrow & L_1F(B) \longrightarrow L_1F(C) \\ & & & & & & \nearrow \delta \\ & & & & F(A) & \longleftarrow & F(B) \longrightarrow F(C) \longrightarrow 0 \end{array}$$

which is natural in the short exact sequence (morphism of δ -functors): if

$$\begin{array}{ccccccc} 0 & \longrightarrow & A & \longrightarrow & B & \longrightarrow & C \longrightarrow 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \longrightarrow & A' & \longrightarrow & B' & \longrightarrow & C' \longrightarrow 0 \end{array}$$

then

$$\begin{array}{ccc} L_{i+1}F(C) & \xrightarrow{\delta} & L_iF(A) \\ \downarrow & & \downarrow \\ L_{i+1}F(C') & \xrightarrow{\delta} & L_{i+1}F(A') \end{array}$$

Remark 1.1. This functor $(L_iF, \delta_i)_i$ satisfies a universal property and is called a δ -functor.

Proof. If we have a short exact sequence of complexes

$$0 \longrightarrow P^A \longrightarrow P^B \longrightarrow P^C \longrightarrow 0$$

this stays exact since $P_i^B = P_i^A \oplus P_i^C$ and F is additive:

$$0 \longrightarrow F(P^A) \longrightarrow F(P^B) \longrightarrow F(P^C) \longrightarrow 0$$

So we get a long exact sequence on cohomology

$$\cdots \longrightarrow H_{i+1}(F(P^C)) \xrightarrow{\delta} H_i(F(P^A)) \longrightarrow H_i(F(P^B)) \longrightarrow H_i(F(P^C)) \longrightarrow \cdots$$

This is

$$\cdots \longrightarrow L_{i+1}F(C) \xrightarrow{\delta} L_iF(A) \longrightarrow L_iF(B) \longrightarrow L_iF(C) \longrightarrow \cdots \quad \square$$

1.2 Injective resolutions and right derived functors

Definition 1.2. A (comological) resolution of an object A is an augmented complex

$$0 \longrightarrow A \longrightarrow C.$$

which is exact. It is **injective** if the objects of C are injective. C has **enough injectives** if for all $A \in \mathcal{C}$, there is an injective $I \in \mathcal{C}$ and a monomorphism $A \rightarrow I$.

Proof. If \mathcal{C} is abelian, then \mathcal{C} has enough injectives if and only if every object has an injective resolution. \square

Proposition 1.2. $R\text{-Mod}$ has enough injectives.

Proof. If $R = \mathbb{Z}$, let A be an abelian group. Then $A \cong \bigoplus_{j \in J} \mathbb{Z}/T_j$, where $T_j \leq \mathbb{Z}$, where J is an index set. This maps into $\bigoplus_{j \in J} \mathbb{Q}/\mathbb{Z}$, which is divisible and hence injective.

If R is arbitrary, let A be an R -module. Let $\phi : \text{Hom}_{\mathbb{Z}}(R, A)$, where $\phi(r) = ra$ is a map of R -modules. As abelian groups, A injects into a divisible abelian group D . So $\text{Hom}_{\mathbb{Z}}(R, A)$ injects into $\text{Hom}_{\mathbb{Z}}(R, D)$. From homework, $\text{Hom}_{\mathbb{Z}}(R, D)$ is injective as an R -module. \square

Definition 1.3. If \mathcal{C} is an abelian category with enough injectives and $F : \mathcal{C} \rightarrow \mathcal{D}$ is left exact, then the i -th **right derived functor** $R^i F : \mathcal{C} \rightarrow \mathcal{D}$ is $R^i F(A) = H^i(F(I_\bullet))$ for $A \rightarrow I_\bullet$ an injective resolution. For $f : A \rightarrow B$ and $g : I_\bullet \rightarrow J_\bullet$ such that

$$\begin{array}{ccc} A & \longrightarrow & I_\bullet \\ f \downarrow & & \downarrow g \\ B & \longrightarrow & J_\bullet \end{array}$$

we have $R^i F(f) = F(g)^*$.

Theorem 1.2. The R^i give a cohomological δ -functor (universal):

$$0 \longrightarrow A \longrightarrow B \longrightarrow C \longrightarrow 0$$

gives a long exact sequence

$$0 \longrightarrow F(A) \longrightarrow F(B) \longrightarrow F(C) \xrightarrow{\delta^0} R^1 F(A) \longrightarrow R^1 F(B) \longrightarrow \dots$$

1.3 Tor functors

Lemma 1.2. Let R be a ring, and let N be an R^{op} -module. Then we have the tensor product functor $t_N : R\text{-Mod} \rightarrow \text{Ab}$ given by $t_N(M) = N \otimes_R M$ and $t_N(f) = \text{id}_N \otimes f$. Then t_N is right exact.

Proof. t_N is additive because direct sums commute with tensor products. If we have a right exact sequence

$$A \xrightarrow{f} B \xrightarrow{g} C \longrightarrow 0$$

then $N \otimes_R C$ is generated by the elements $n \otimes c$ with $c \in C$ and $n \in N$. If $c = g(b)$, then $t_N(g)(b) = n \otimes c$, so $t_N(g)$ is surjective. So we get a surjection $\text{coker}(t_N(f)) \rightarrow N \otimes_R C$. We want this to be an isomorphism, so let's try to create an inverse. Define an R -balanced map $N \times C \rightarrow \text{coker}(T_N(f))$ by $\theta(n, c) = n \otimes b + \text{im}(t_n(f))$. Then $g(b) = c$. This is well-defined because if $g(b') = c$, then $g(b - b') = 0$, so $b - b' = f(a)$ for some a ; then $n \otimes b - n \otimes b' + t_N(f)(a)$. So we've constructed an inverse to the surjection. \square

Definition 1.4. The i -th Tor functor is $\text{Tor}_i^R(N, \cdot) := L_i t_N$.